

# CS250P: Computer Systems Architecture

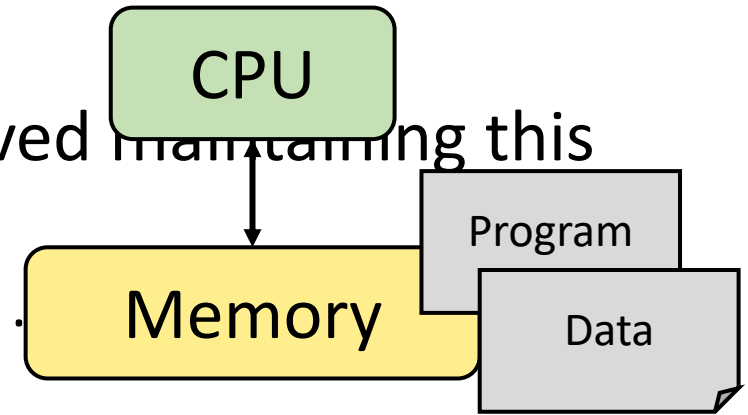
## Moore's Law



Sang-Woo Jun  
Fall 2022

# Conventional performance scaling

- ❑ Traditional model of a computer is simple
  - Single, in-order flow of instructions on a processor
  - Simple, in-order memory model
- ❑ Large part of computer architecture research involved maintaining this abstraction while improving performance
  - Transparent caches, Transparent superscalar scheduling, .
  - Same software runs faster tomorrow
  - (Slow software becomes acceptable tomorrow)
- ❑ Driven largely by continuing march of Moore's law



# Moore's Law

- ❑ What exactly does it mean?
- ❑ What is it that is scaling?

# Moore's Law

- ❑ Typically cast as:

“Performance doubles every X months”

- ❑ Actually closer to:

“Number of transistors per unit cost doubles every X months”

# Moore's Law

The complexity for minimum component costs has increased at a rate of roughly a factor of two per year.

[...]

Over the longer term, the rate of increase is a bit more uncertain, although there is no reason to believe it will not remain nearly constant for at least 10 years.

-- Gordon Moore, Electronics, 1965

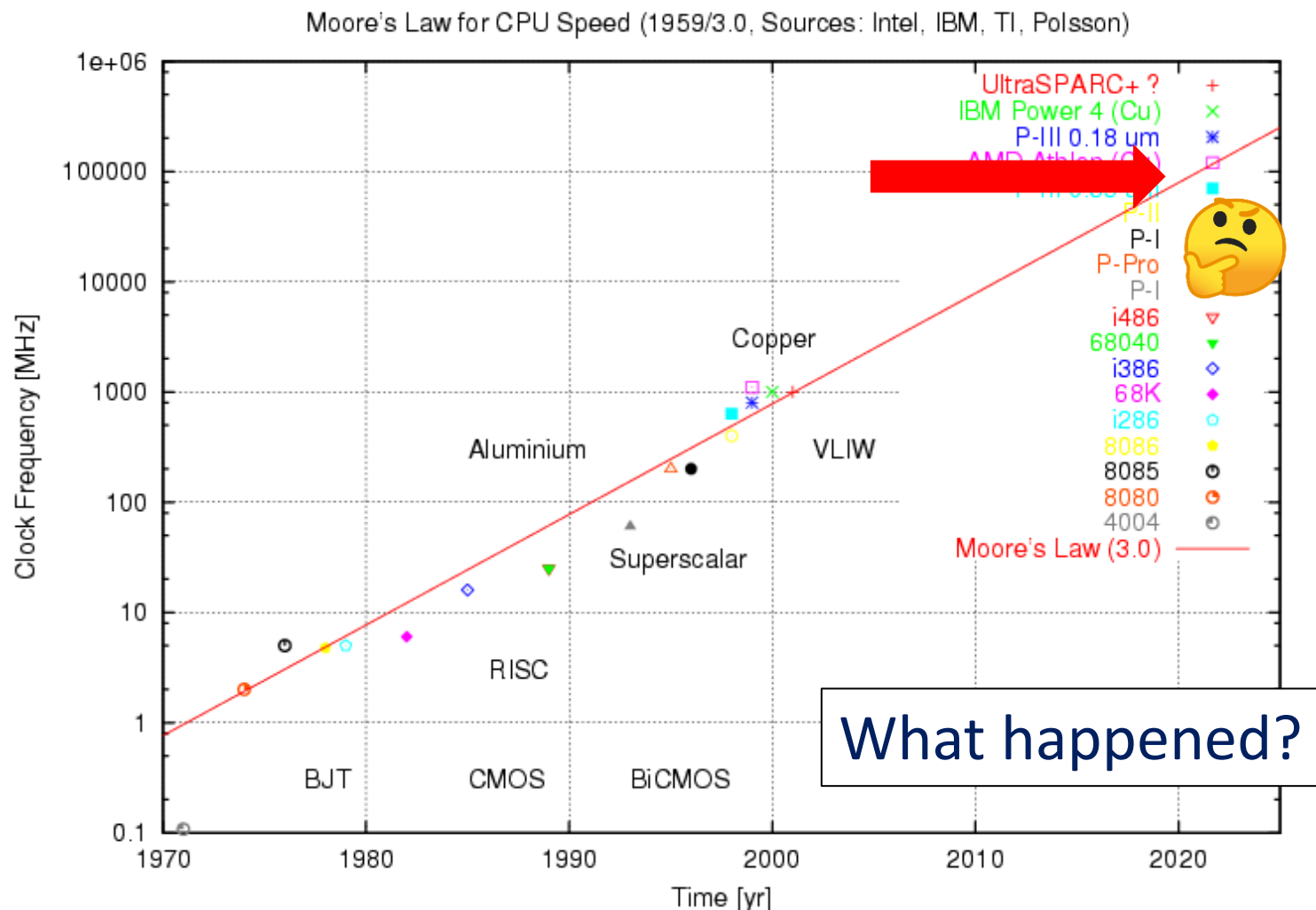
Why is Moore's Law conflated with processor performance?

# Dennard Scaling: From Moore's Law to performance

- ❑ “Power density stays constant as transistors get smaller”
  - Robert H. Dennard, 1974
  
- ❑ Intuitively:
  - Smaller transistors → shorter propagation delay → faster frequency
  - Smaller transistors → smaller capacitance → lower voltage
  
  - $Power \propto Capacitance \times Voltage^2 \times Frequency$

Moore's law → Faster performance @ Constant power!

# Single-core performance scaling projection



# (Slightly) more accurate processor power consumption

Gate-oxide    Stopped scaling  
stopped scaling    due to leakage

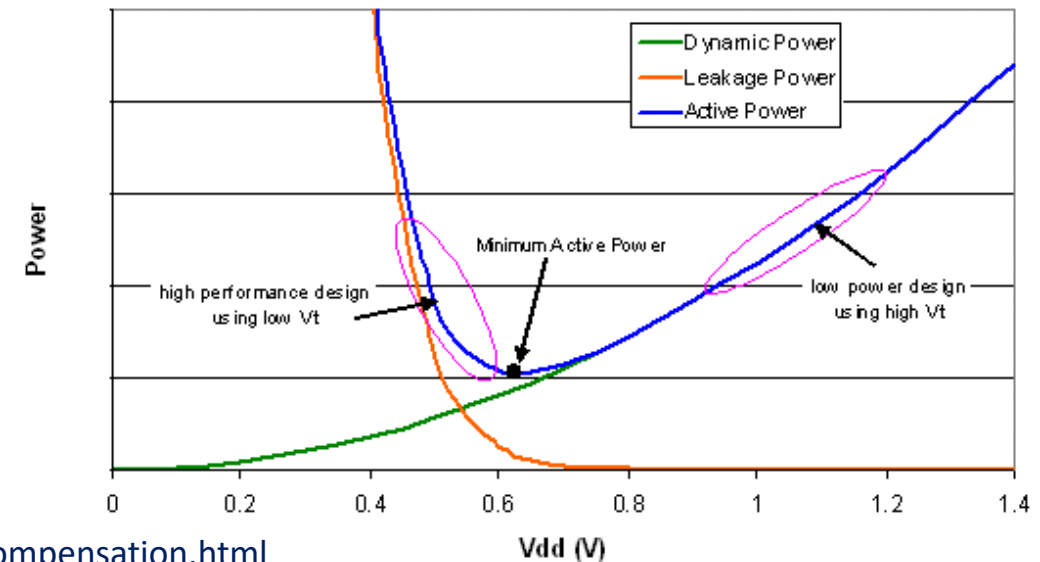
$$Power = \underbrace{(ActiveTransistors \times Capacitance \times Voltage^2 \times Frequency)}_{\text{Dynamic power}} + \underbrace{(Voltage \times Leakage)}_{\text{Static power}}$$

Dynamic power

$$+ (Voltage \times Leakage)$$

Static power

Total power consumption with constant frequency

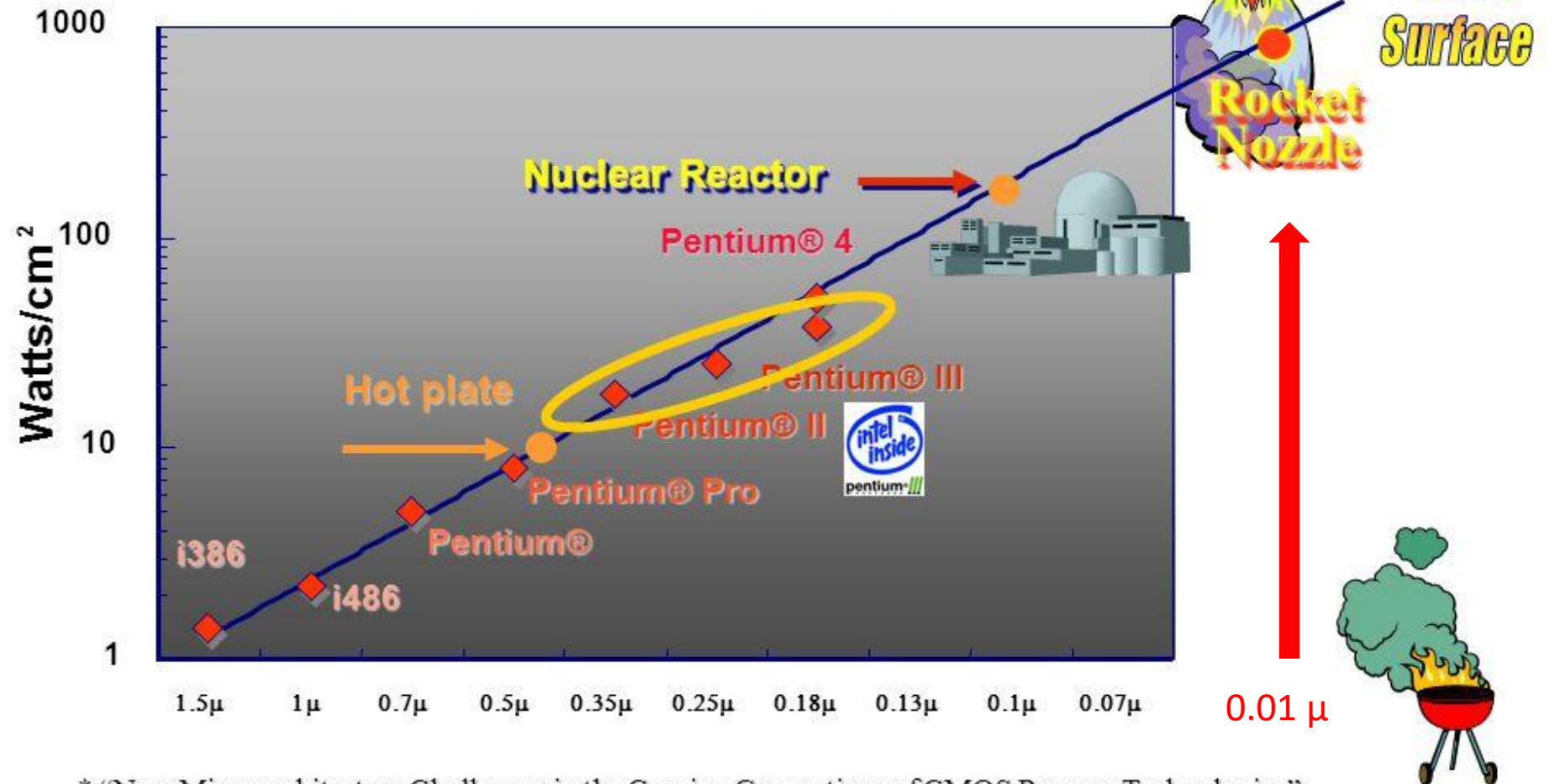




# End of Dennard Scaling

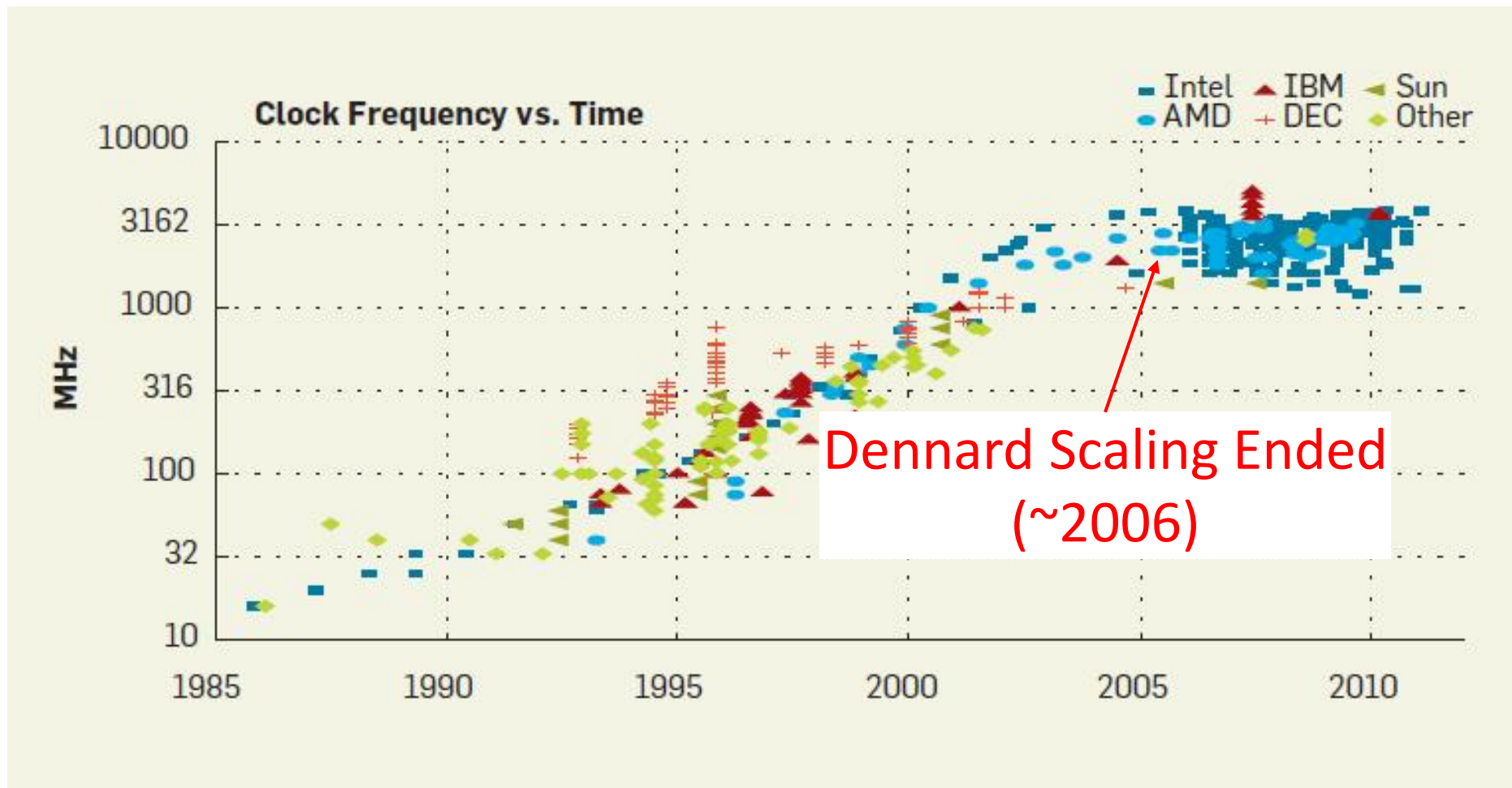
- ❑ Even with smaller transistors, we cannot continue reducing power
  - What do we do now?
  
- ❑ Option 1: Continue scaling frequency at increased power budget
  - Chip quickly become too hot to cool!
  - Thermal runaway:
    - Hotter chip → increased resistance → hotter chip → ...

# Option 1: Continue scaling frequency at increased power budget

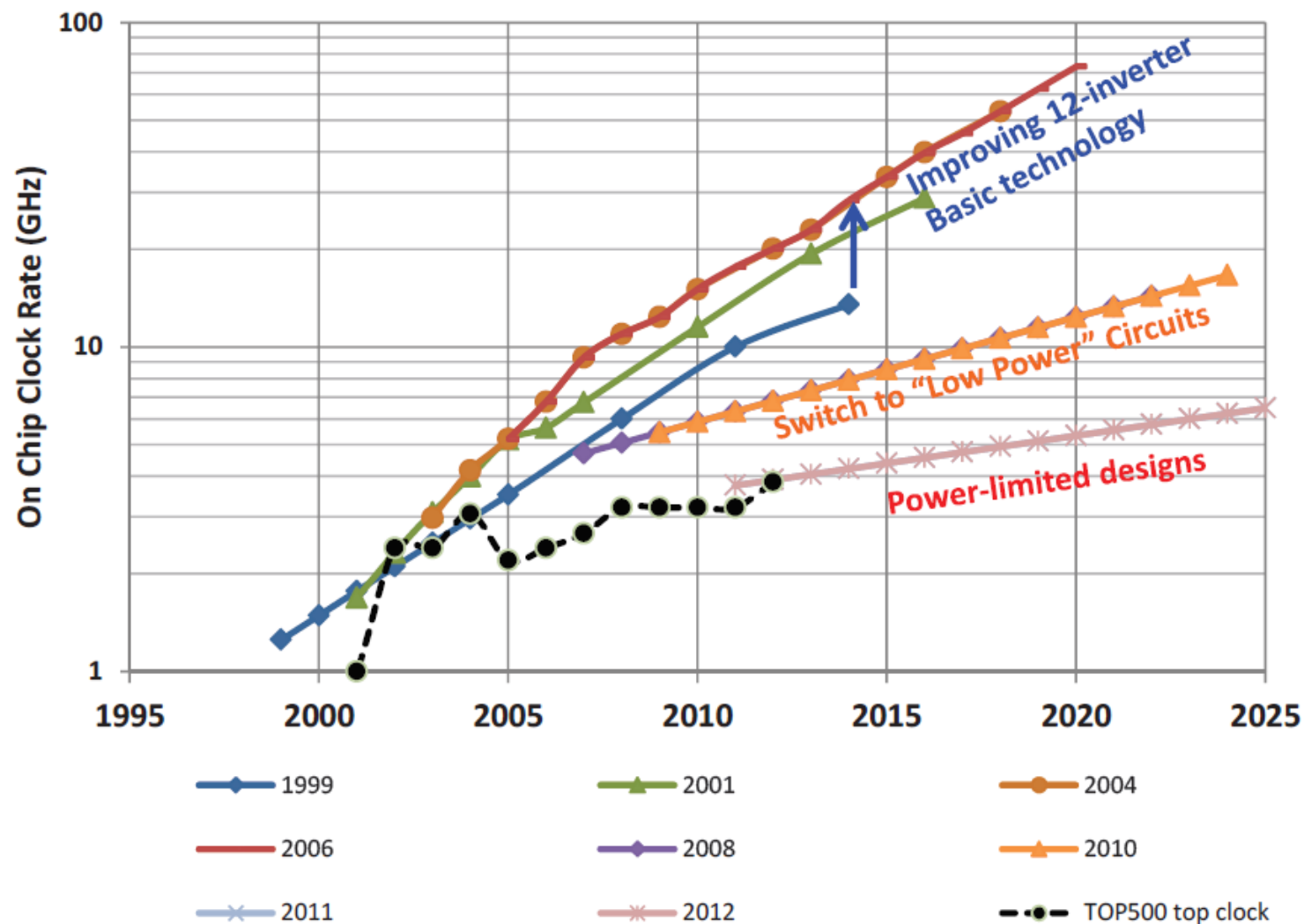


\* "New Microarchitecture Challenges in the Coming Generations of CMOS Process Technologies" – Fred Pollack, Intel Corp. Micro32 conference key note - 1999.

# Option 2: Stop frequency scaling



# Looking back: change of predictions

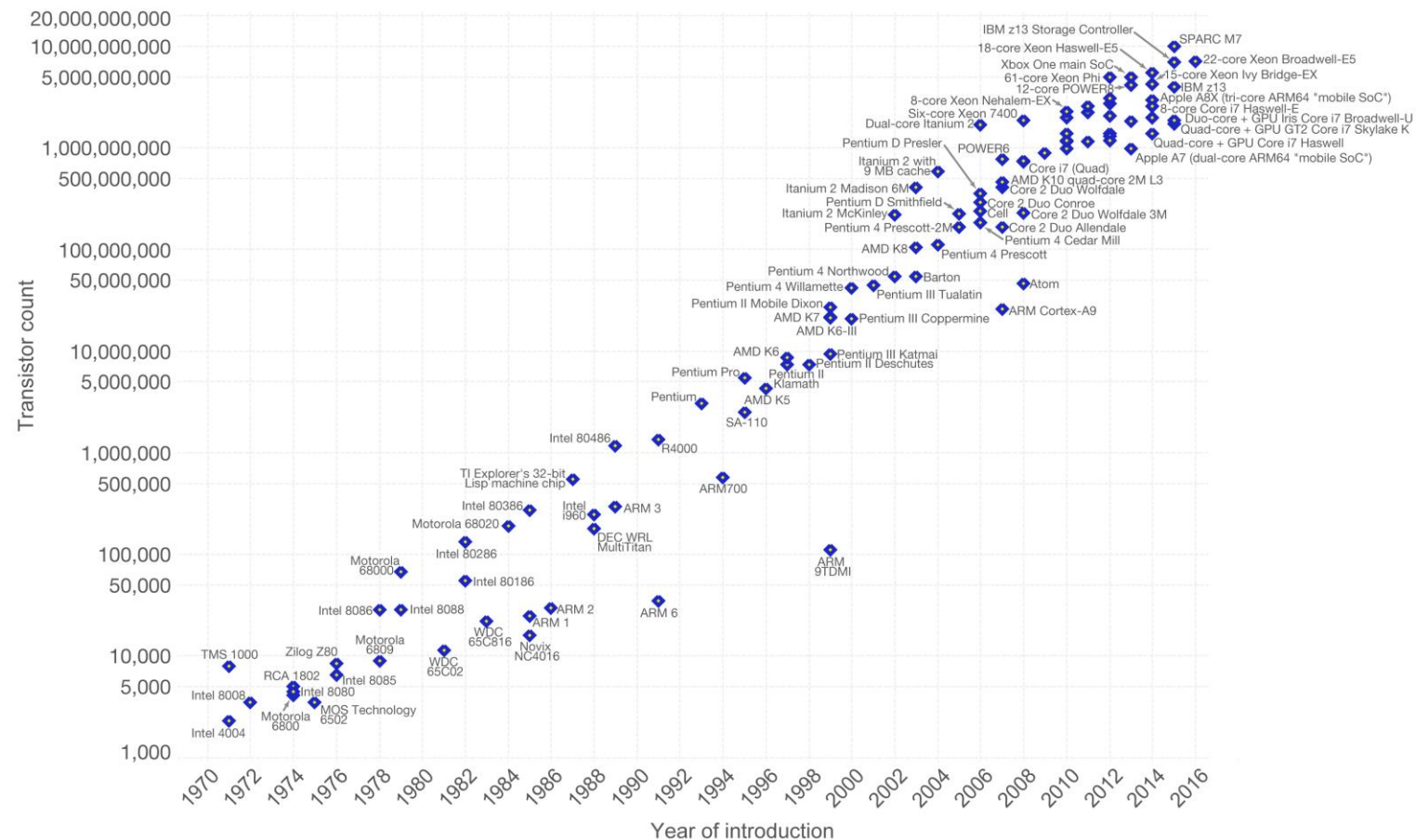


# But Moore's Law continues beyond 2006

## Moore's Law – The number of transistors on integrated circuit chips (1971-2016)



Moore's law describes the empirical regularity that the number of transistors on integrated circuits doubles approximately every two years. This advancement is important as other aspects of technological progress – such as processing speed or the price of electronic products – are strongly linked to Moore's law.



Data source: Wikipedia ([https://en.wikipedia.org/wiki/Transistor\\_count](https://en.wikipedia.org/wiki/Transistor_count))

The data visualization is available at [OurWorldinData.org](https://www.ourworldindata.org). There you find more visualizations and research on this topic.

Licensed under CC-BY-SA by the author Max Roser.

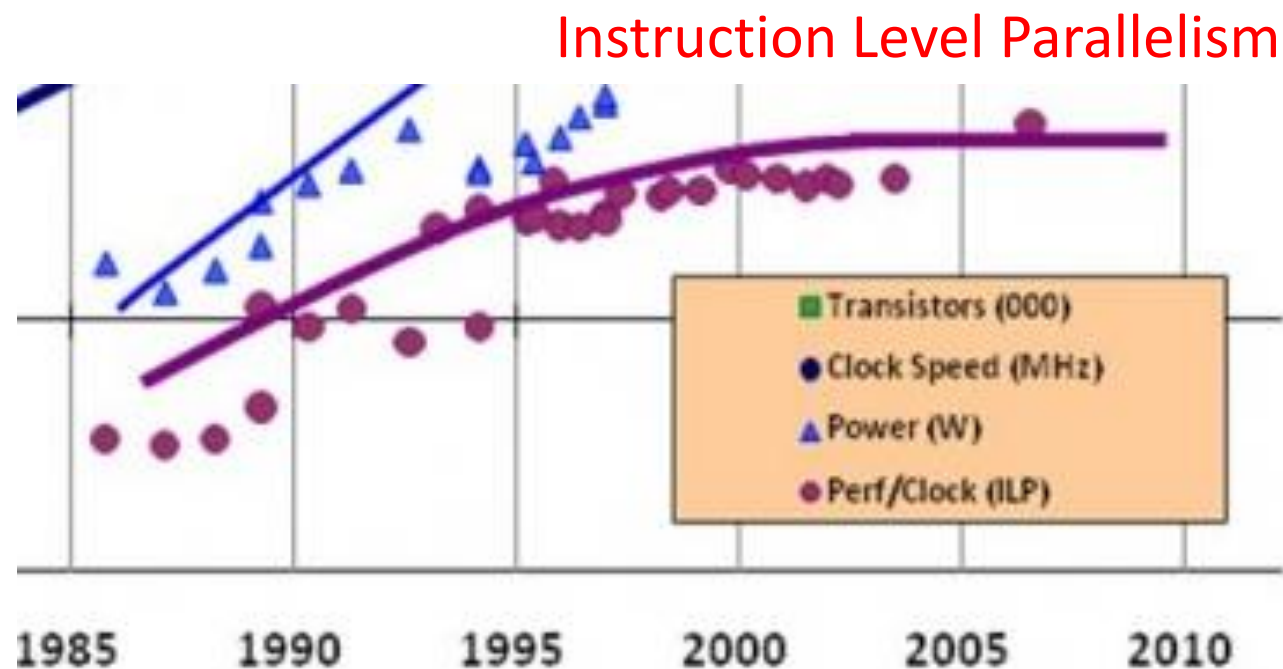
# State of things at this point (2006)

## ❑ Single-thread performance scaling ended

- Frequency scaling ended (Dennard Scaling)
- Instruction-level parallelism scaling stalled ... also around 2005

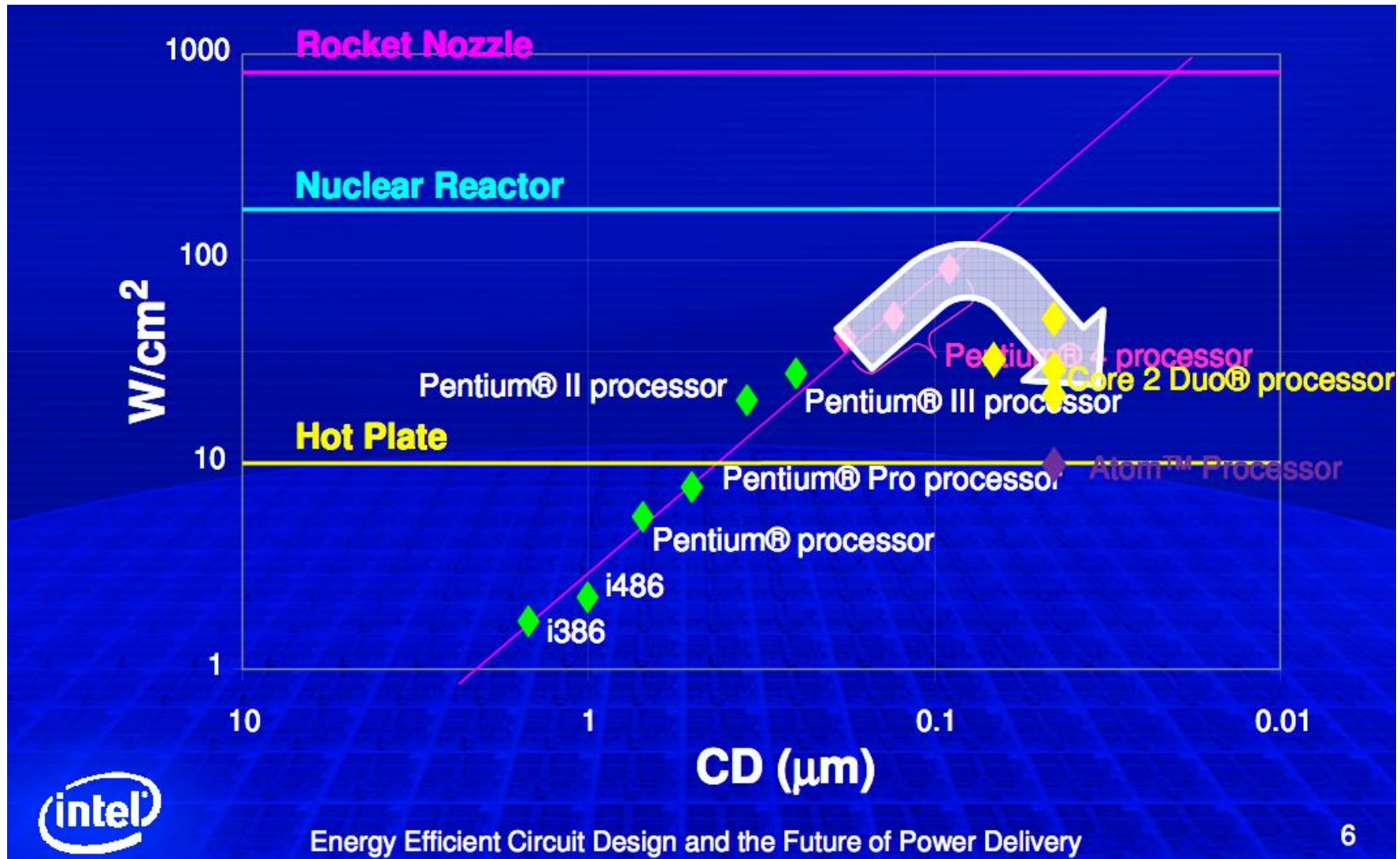
## ❑ Moore's law continues

- Double transistors every two years
- What do we do with them?

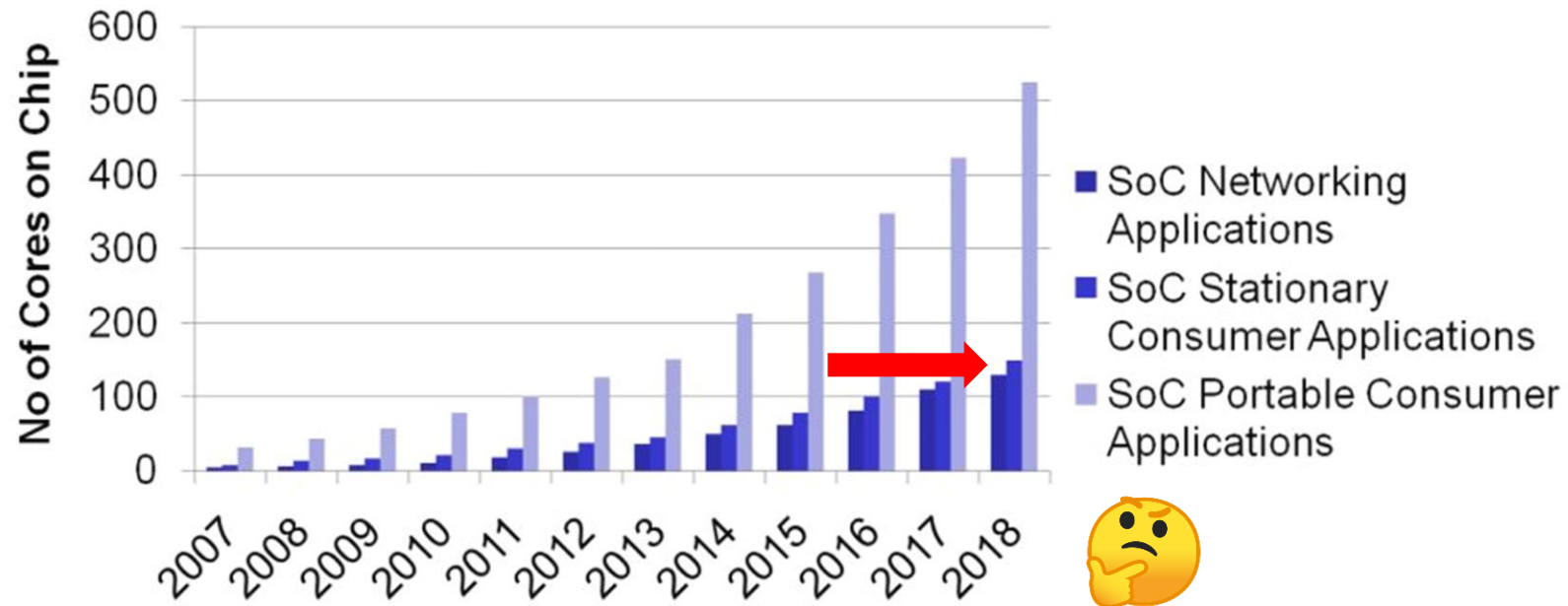




# Crisis averted with manycores?



# Crisis averted with manycores?



Source:

International Roadmap for Semiconductors 2007 edition (<http://www.itrs.net/>)



# What happened?

Can't keep going up

$$\text{Power} = (\text{ActiveTransistors} \times \text{Capacitance} \times \text{Voltage}^2 \times \text{Frequency}) + (\text{Voltage} \times \text{LeakageCurrent})$$

Gate-oxide stopped scaling    Stopped scaling due to leakage    Stopped scaling due to thermal

Dynamic power

Static power

“Utilization Wall”

Regardless of Moore's Law, a limited amount of gates can be active at a given time

# Where To, From Here?

- ❑ The number of active transistors at a given time is limited
  - We won't get much performance improvements even if Moore's law continues
  - We need to make the best use of those active transistors!

# Where To, From Here?

## ❑ Potential Solution 1: The software solution

- Write efficient software to make the efficient use of hardware resources
- No longer depend entirely on hardware performance scaling
- “Performance engineering” software, using hardware knowledge



## ❑ Solution 2: The specialized architectural solution

- Chip space is now cheap, but power is expensive
- Stop depending on more complex general-purpose cores
- Use space to build heterogeneous systems, with compute engines well-suited for each application



# The Bottom Line:

## Architecture is No Longer Transparent

- ❑ Optimized software requires architecture knowledge
- ❑ Special-purpose “accelerators” (GPU, FPGA, ...) programmed explicitly
- ❑ Even general-purpose processors implement specialized instructions
  - Single-Instruction Multiple Data (SIMD) instructions such as AVX
  - Special-purpose instructions sets such as AES-NI